



## DIGITAL CIRCUITS SOLUTIONS TO EXAMPLARY EXAM QUESTIONS

### SOLUTION 1:

a. B is negative, result is negative, there is an overflow, and operation is subtraction

i) Overflow condition: pos – neg = neg , therefore A must be **positive**.

$$\begin{array}{rcl} \text{ii) } A = 0xxx \ xxxx & 0xxx \ xxxx & \\ B = 1001 \ 1101 & 2's \ comp. \ + \ 0110 \ 0011 & \text{smallest possible } A = 0001 \ 1101 \\ R = 1xxx \ xxxx & 1xxx \ xxxx & \end{array}$$

The same solution by thinking in decimal:

$B = (-99)_{10}$  , to generate an overflow result must be at least +128. (Note that result seems to be negative, but due to overflow the real sign of the result is positive.)

$A - 99 = 128$ , smallest possible  $A = (29)_{10} = \mathbf{0001 \ 1101}$

b. The carry bit is 1. It means **no borrow**. Therefore  $A > B$ .

### SOLUTION 2:

a)

Expression:

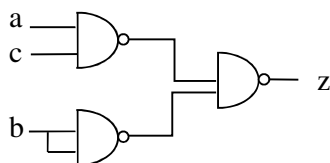
$$(a+E)(a'+F) \text{ or } (a'+E)(a+F)$$

b)

$$\begin{aligned} (a+E)(a'+F)(E+F) &= (a+E)(a'+F)(E+F+aa') && \text{Inverse and identity} \\ &= (a+E)(a'+F)(E+F+a)(E+F+a') && \text{Distribution} \\ &= (a+E)(1+F)(a'+F)(1+E) && \text{Identity} \\ &= (a+E)(a'+F) \end{aligned}$$

c)

$$\begin{aligned} z &= ab'c + acd + ab + a'b \\ &= ab'c + acd + b(a'+a) && \text{Distribution} \\ &= acd + ab'c + b && \text{Inverse} \\ &= acd + ab'c + b + ac && \text{Consensus} \\ &= ac(d+b'+1) + b && \text{Distribution and identity} \\ &= ac + b \end{aligned}$$



### SOLUTION 3:

a)

$$\begin{aligned}
 f(A,B,C,D) &= A'B'CD + AB'CD + AC'D + AC'D' + A'B'CD' + ABCD + ACD' \\
 &= (A'+A)B'CD + AC'(D+D') + A'B'CD' + ABCD + ACD' \quad (\text{Inverse}) \\
 &= (B'+AB)CD + AC' + A'B'CD' + ACD' \quad (\text{absorbition}) \\
 &= B'CD + A(CD + C') + (A'B'+A)CD' \quad (\text{absorbition}) \\
 &= B'CD + AD + AC' + B'CD' + ACD' \\
 &= B'CD + AD + B'CD' + A(C'+CD') \quad (\text{absorbition}) \\
 &= B'CD + AD + B'CD' + AC' + AD' \\
 &= B'C(D + D') + A(D+C'+D') \quad (\text{inverse}) \\
 &= B'C + A
 \end{aligned}$$

b)

f		CD			
		00	01	11	10
AB	00	1	1	1	1
	01	0	0	0	1
	11	Φ	0	Φ	1
	10	1	1	1	1

By considering 0 and Φ points we can obtain complement of f.

$$\bar{f} = B\bar{C} + BD$$

De Morgan:

$$\begin{aligned}
 f &= \overline{B\bar{C} + BD} = \overline{B\bar{C}} \cdot \overline{BD} \\
 &= (\bar{B} + C) \cdot (\bar{B} + \bar{D})
 \end{aligned}$$

Or by considering true (1) points:

$$f = \bar{B} + (C\bar{D})$$

Distributive Law:

$$f = (\bar{B} + C) \cdot (\bar{B} + \bar{D})$$

### SOLUTION 4 :

a) Maxterms (0 generating inputs): 0001, 0101, 1100, 1101, 1110, 1001

cd		ab			
		00	01	11	10
ab	00	1		1	1
	01	1		1	1
	11			1	
	10	1		1	1

Prime Implicants:

b'd' , a'd' , a'c, b'c, cd

b)

False (0) points of function f are true (1) points of the **complement** ( $\bar{f}(a,b,c,d)$ ).

Num	abcd	Num	abcd	Num	abcd
1	0001√	1,5	0-01√	1,5,9,13	- - 01 <b>X</b>
5	0101√	1,9	-001√		
9	1001√	5,13	-101√		
12	1100√	9,13	1-01√		
13	1101√	12,13	110- <b>X</b>		
14	1110√	12,14	11-0 <b>X</b>		

Prime Implicants:

abc', abd', c'd

### SOLUTION 5 :

**Step 1.** Point 4 is a distinguished point, and F is an essential prime implicant. F is selected, points 4, 7, 13 and 14 are removed.

**Step 2.** C covers D with equal cost. D is removed.

**Step 3.** True point 2 is a distinguished point, and C is an essential prime implicant. C is selected, points 2, 5, 8 and 10 are removed.

**Step 4.** B covers E with less cost. E is removed.

**Step 5.** True point 11 is a distinguished point, and B is an essential prime implicant. B is selected.

Hence B+C+F is the minimal covering sum (sufficient base) with 26 unit cost.

### SOLUTION 6:

a)

**Z**

CD \ AB	00	01	11	10
00	1	0	1	1
01	1	1	1	0
11	Φ	0	0	0
10	1	0	1	1

Set of all prime implicants:

$\bar{C}\bar{D}$ ,  $\bar{B}C$ ,  $\bar{B}\bar{D}$ ,  $\bar{A}B\bar{C}$ ,  $\bar{A}BD$ ,  $\bar{A}CD$   
A B C D E F

b)

### Prime Implicant Chart:

	0	2	3	4	5	7	8	10	11	Cost
✓ A	X			X			X			6
✓ B		X	X					X	(X)	5
C	X	X					X	X		6
D				X	X					8
✓ E					X	X				7
F			X			X				7

c) Cheapest sufficient set of prime implicants:

A + B + E: Cost=18

Cheapest expression:  $Z = \bar{C}\bar{D} + \bar{B}C + \bar{A}BD$

### SOLUTION 7:

a.

Truth table:

a	b	c	d	z
0	0	0	0	0
0	0	0	1	0
0	0	1	0	1
0	0	1	1	1
0	1	0	0	1
0	1	0	1	0
0	1	1	0	1
0	1	1	1	0
1	0	0	0	1
1	0	0	1	1
1	0	1	0	1
1	0	1	1	1
1	1	0	0	1
1	1	0	1	1
1	1	1	0	0
1	1	1	1	0

From the truth table we can obtain minterms and write the expression of the function in the 1<sup>st</sup> canonical form.

Remember, in minterms all variables (literals) appear once.

1<sup>st</sup> canonical form:

$$z = a'b'cd' + a'b'cd + a'bc'd' + a'bcd' + ab'c'd' + ab'cd + ab'cd' + ab'cd + abc'd' + abc'd$$

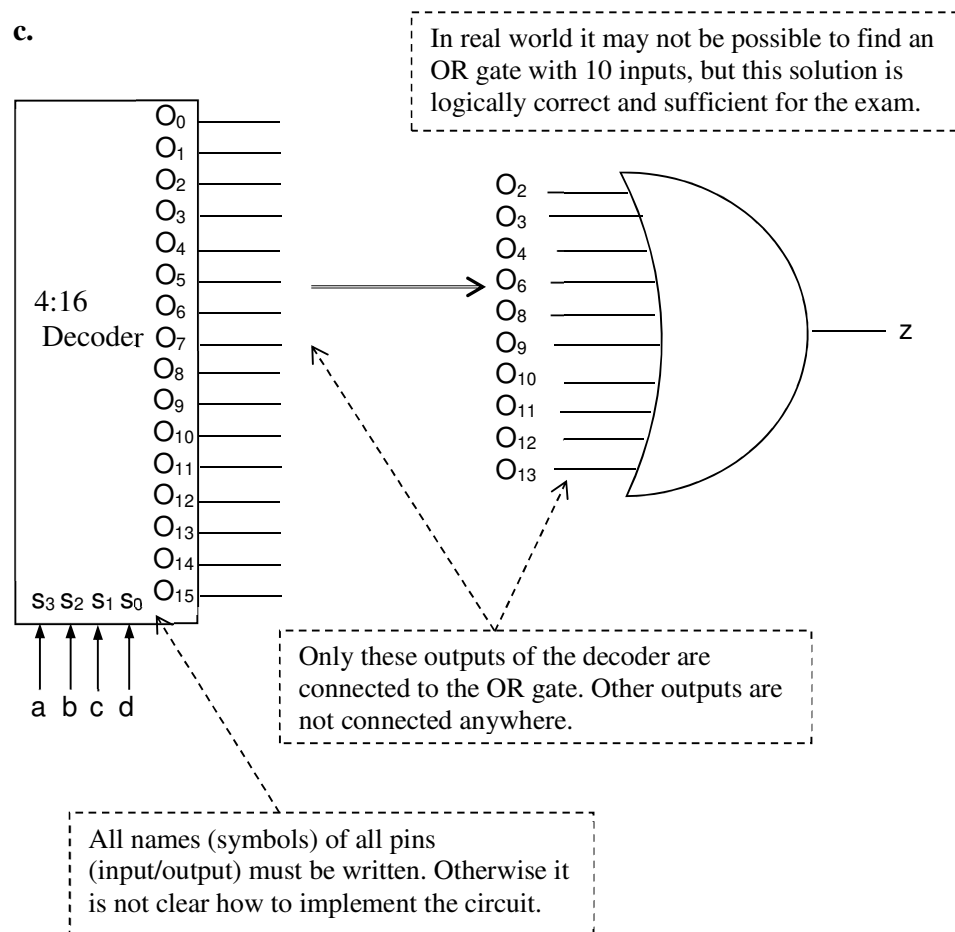
b.

There are different ways to minimize the expression in the 1<sup>st</sup> canonical form.

Minimized expression:

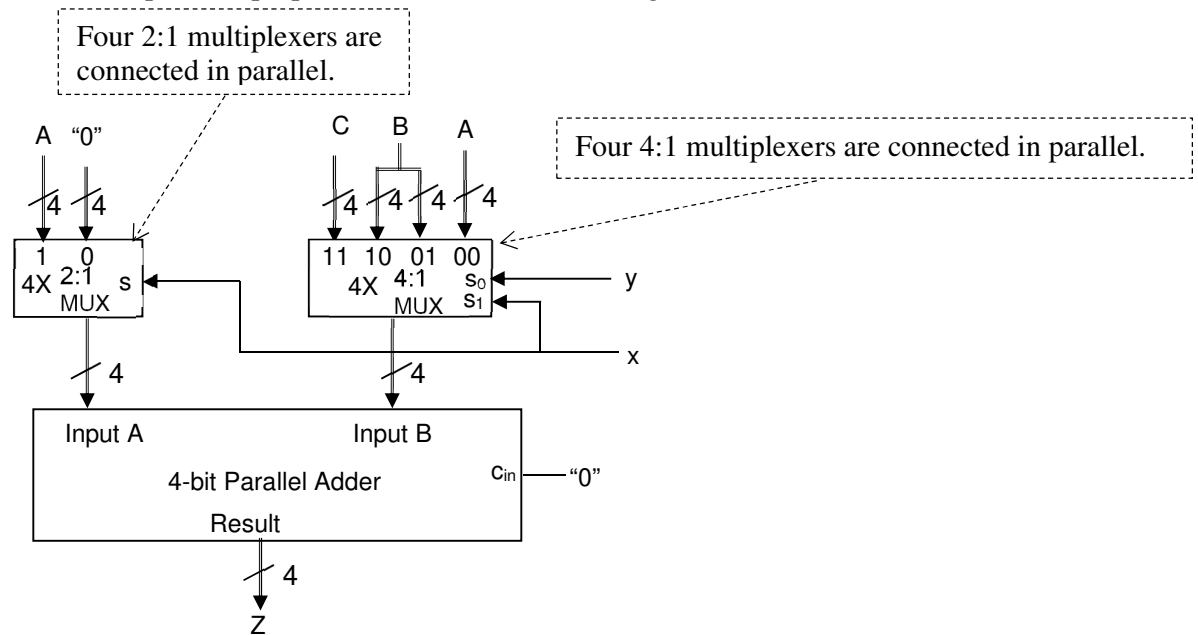
$$z = a'bd' + ac' + b'c$$

c.

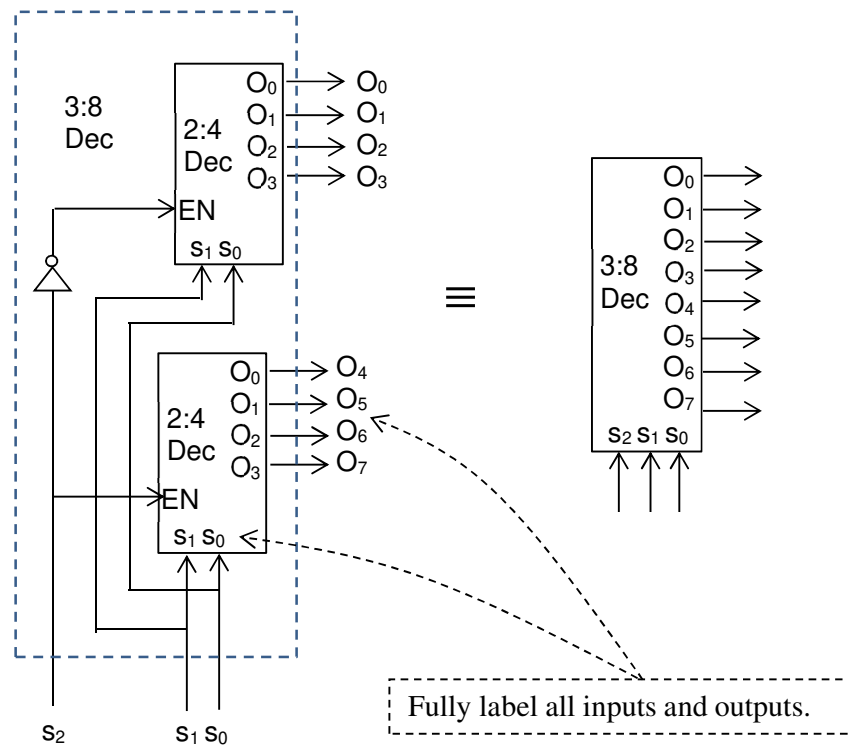


### SOLUTION 8:

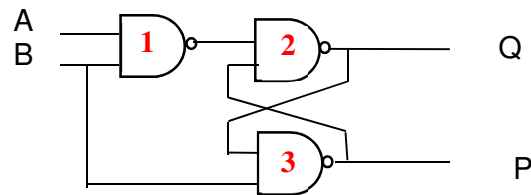
There are different possible proper solutions. One of them is given below.



### SOLUTION 9:



# **SOLUTION 10:**



A	B	Q	P	
0	0	0	1	If B=0 then P = 1. The output of NAND1 is 1 and Q is 0 (stable)
0	1	0	1	After (A=0, B=0), Q=0 from previous state and P = 1. The output of NAND1 is 1 and Q = 0 (stable)
1	0	0	1	If B=0 then P = 1. The output of NAND1 is 1 and Q is 0 (stable)
1	1	1	0	The output of NAND1 is 0. Therefore the output of NAND2, Q= 1 and P=0 (stable)
0	1	1	0	After (A=1, B=1), The output of NAND1 is 1. P = 0 from previous state and Q = 1, P=0 (stable)

A	B	
0	0	Reset
1	0	
0	1	Don't change
1	1	Set

The circuit is stable and can be switched to another state and has set, reset and don't change conditions. Therefore it can be used as a memory unit.

A	B	Q(t)	Q(t+1)
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1

AB\Q(t)	0	1
00	0	0
01	0	1
11	1	1
10	0	0

$Q(t+1) = BQ(t) + AB$

# **SOLUTION 11:**

$$a) Q_1^+ = A\bar{Q}_1 + \bar{B}Q_1$$

$$Q_0^+ = 1 \oplus Q_0 = \bar{Q}_0$$

$$Z = \bar{S}_1 S_0 + S_1 \bar{S}_0$$

$$= S_1 \oplus S_0 = Q_1 \oplus Q_0$$

Output is a function of states  $\rightarrow$  Moore Model

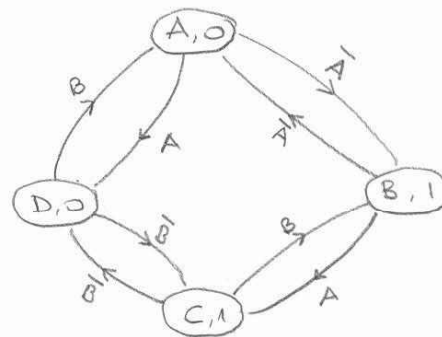
$Q_1^+ Q_0^+$	$AB$				$Z$
$Q_1 Q_0$	00	01	10	11	
00	01	01	11	11	0
01	00	00	10	10	1
10	11	01	11	01	1
11	10	00	10	00	0

A: 00

B: 01

C: 10

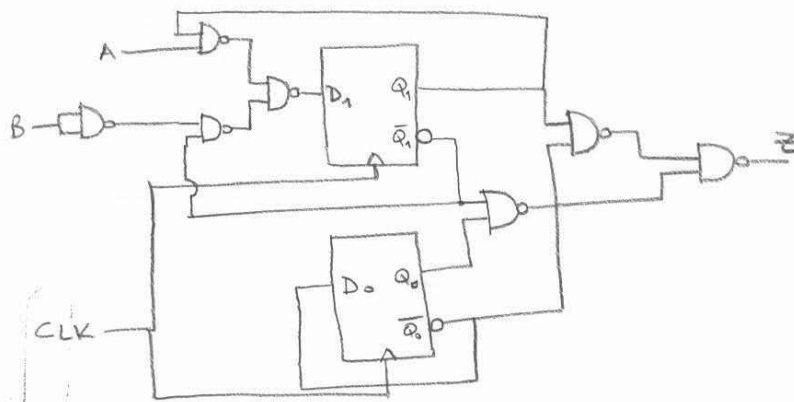
D: 11



$$b) Q_1^+ = D_1 = A\bar{Q}_1 + \bar{B}Q_1$$

$$Q_0^+ = D_0 = \bar{Q}_0$$

$$Z = Q_1 \oplus Q_0$$



# **SOLUTION 12:**

a)

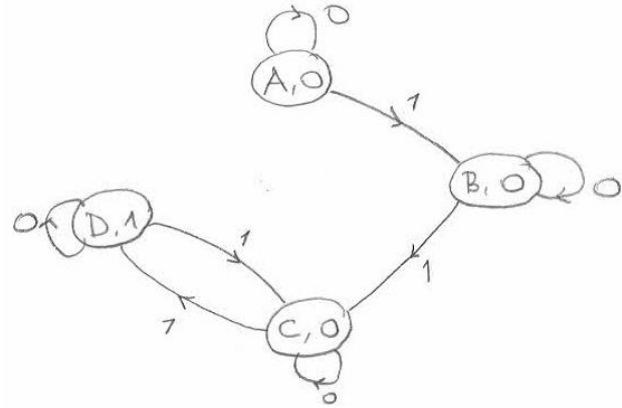
States

A: zero ones

B: one ones

C: two or even ones

D: (three or more) and odd ones



State/output table

$Q_1^+ Q_0^+, Z$

$Q_1 Q_0$	$X$		$Z$
	0	1	
A	A	B	0
B	B	C	0
C	C	D	0
D	D	C	1

Codes

A: 00  
B: 01  
C: 10  
D: 11

$Q_1^+ Q_0^+$

$Q_1 Q_0$	$X$		
	0	1	
00	00	01	
01	01	10	
10	10	11	
11	11	10	

b)

Symbol	Transition	J	K
0	0 → 0	0	φ
α	0 → 1	1	φ
β	1 → 0	φ	1
1	1 → 1	φ	0

$Q_1^+$	$X$		
$Q_1 Q_0$	0	1	
00	0	0	
01	0	α	
10	1	1	
11	1	1	

$Q_0^+$	$X$		
$Q_1 Q_0$	0	1	
00	0	α	
01	1	β	
10	0	α	
11	1	β	



$J_1$	$X$	
$Q_1 Q_0$	0	1
00	0	0
01	0	1
11	$\phi$	$\phi$
10	$\phi$	$\phi$

$K_1$	$X$	
$Q_1 Q_0$	0	1
00	$\phi$	$\phi$
01	$\phi$	$\phi$
11	0	0
10	0	0

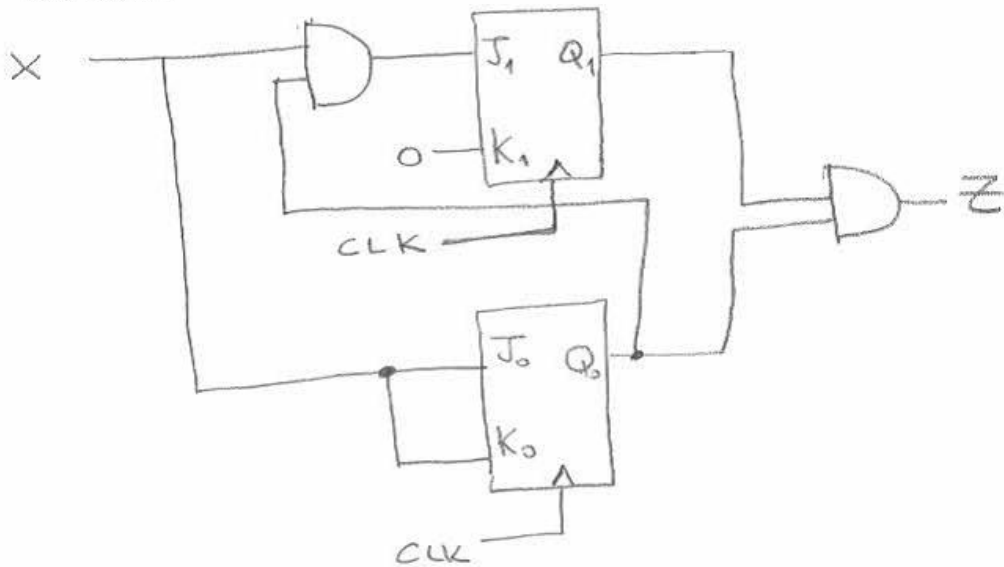
$J_0$	$X$	
$Q_1 Q_0$	0	1
00	0	1
01	$\phi$	$\phi$
11	0	1
10	$\phi$	$\phi$

$K_0$	$X$	
$Q_1 Q_0$	0	1
00	$\phi$	$\phi$
01	0	1
11	$\phi$	$\phi$
10	0	1

$$J_1 = Q_0 X \quad J_0 = X$$

$$K_1 = 0 \quad K_0 = X$$

$$Z = Q_1 Q_0$$



### SOLUTION 13:

A	Q1	Q2	Q3	Q4	Z
L	Off	On	On	Off	L
H	On	Off	Off	On	H

The expression for the function  $Z = f(A) = A$  (buffer).